COMP 212 Spring 2025 Homework 07

In lecture and lab, we looked at higher-order functions on lists. Since lists are bad for parallelism, so in this homework we will investigate similar higher-order functions on trees.

Previously, we used trees that had data at each internal node. For this week, we will instead consider trees where there is data only at the leaves:

An 'a tree is either

- 1. empty
- 2. a leaf with value x:'a
- 3. a node with two subtrees

and that's it!

This is represented by the following datatype:

```
datatype 'a tree = Empty | Leaf of 'a | Node of 'a tree * 'a tree
```

This datatype is defined in lib.sml, which is loaded by your homework file. lib.sml also provides

```
val fromlist : 'a list -> 'a tree
val tolist : 'a tree -> 'a list
```

which you can use for testing (but only for testing!). Since we are mostly concerned with the contents of the trees, and not their specific arrangement, we can write tests that convert lists to trees and back, which has a more concise notation than writing the trees themselves. There are some examples in the homework code, which you can add more tests to.

We have also provided a function sort

```
sort : ('a * 'a -> order) * 'a tree -> 'a tree
```

based on the tree sorting code from earlier in the semester (together with some additional code to convert the leaf-empty-node trees used here to the leaf-node trees used in sorting, and vice versa).

sort takes a comparison function, which produces an order, which is datatype whose constructors are LESS and EQUAL and GREATER:

```
datatype order = LESS | EQUAL | GREATER
```

It sorts the tree into increasing order according to this comparison function. There are built-in comparison functions on integers (Int.compare) and floating points (Real.compare). For example,

1 Map

Consider the following two functions:

```
(* pluralize_rec t evaluates to a tree t', where
  t' has the same structure as t, and the string at each leaf of t'
  is the string at the corresponding position in t, with an 's'
  affixed to the end. *)
fun pluralize_rec (t : string tree) : string tree =
    case t of
        Empty => Empty
      | Leaf x \Rightarrow Leaf (x ^ "s")
      | Node(l,r) => Node(pluralize_rec l , pluralize_rec r)
(* mult_rec t evaluates to a tree t', where
   t' has the same structure as t, and the int at each leaf of t'
   is the int at the corresponding position in t, multiplied by c
*)
fun mult_rec (c : int, t : int tree) : int tree =
    case t of
        Empty => Empty
      | Leaf x \Rightarrow Leaf (c * x)
      | Node(l,r) => Node(mult_rec (c,l) , mult_rec (c,r))
```

The pattern is "compute a new tree by applying some function **f** to each element of the given tree."

Task 1.1 (8 pts). Write a higher-order function

that abstracts this pattern.

Task 1.2 (2 pts). Rewrite pluralize and mult using map. The functions should be equivalent to pluralize_rec and mult_rec, but should be defined using map rather than being defined directly by recursion.

2 Reduce

Consider the following two functions:

```
(* sum_rec t evaluates to a number n, where n is
   the sum of all of the numbers at the leaves of t *)
fun sum_rec (t : int tree) : int =
    case t of
        Empty => 0
      | Leaf x \Rightarrow x
      | Node(t1,t2) => (sum_rec t1) + (sum_rec t2)
(* join_rec t evaluates to a string s, where s is
   the concatenation of all of the strings at the leaves of t,
   in order from left to right *)
fun join_rec (t : string tree) : string =
    case t of
        Empty => ""
      | Leaf x \Rightarrow x
      | Node(t1,t2) => (join_rec t1) ^ (join_rec t2)
val "abc" = join_rec(Node(Leaf "a", Node(Leaf "b", Leaf "c")))
```

The general pattern here is called reduce, which takes a binary function of type 'a * 'a -> 'a to apply at each node, and a value of type 'a for the empty tree, and computes an 'a from an 'a tree.

Task 2.1 (8 pts). Write the function

```
reduce : ('a * 'a -> 'a) * 'a * 'a tree -> 'a
```

that implements the operation of reduction on trees.

Task 2.2 (2 pts). Rewrite sum and join using reduce. The functions should be equivalent to sum_rec and join_rec, but should be defined using reduce rather than being defined directly by recursion.

3 Programming with map and reduce

To receive credit for a task in this section, your function must not be defined recursively. The goal is to practice programming by combining higher-order functions. In each task, you may use map, reduce, any previous tasks in this section, and any other functions that are specifically allowed by the task. If you cannot figure out how to solve a task this way, you may wish to first define it recursively, and then think about how the recursive version can be rewritten with higher-order functions.

We say that x is an *element* of a tree t if Leaf x occurs somewhere in t. In all of these tasks, the shape of the resulting tree is up to you, as long as it has the correct elements in the order specified in the problem. For example, the following trees have all of the same elements in the same order:

```
Node(t1,Node(t2,t3)) and Node(Node(t1,t2),t3)
Node(Empty,t) and t
Node(t,Empty) and t
```

so one side is always just as correct an answer as the other. In particular, you never need to rebalance a tree.

3.1 Flatten

The type ('a tree) tree represents a tree whose leaves are themselves trees.

```
Task 3.1 (5 pts). Write a function
```

```
flatten : ('a tree) tree -> 'a tree
```

such that flatten t contains all of the elements of all of the trees in t. The elements of each tree t1 in t should occur in flatten t in the same order in which they occur in t1; if a tree t1 is to the left of a tree t2 in t, the elements of t1 should occur to the left of the elements of t2 in flatten t.

For example:

3.2 Filter

```
Task 3.2 (5 pts). Define a function
```

```
filter: ('a -> bool) * 'a tree -> 'a tree
```

such that filter (p, t) contains all and only the elements x : 'a of t for which p x returns true. The elements that are kept should be in the same order as in the original tree.

For example:

```
filter (fn x => x > 2, Node (Node (Leaf 1,Leaf 2),Node (Leaf 3,Empty))) == Node (Node (Empty,Empty),Node (Leaf 3,Empty))
```

Hint: first create an ('a tree) tree where each Leaf x in the original tree is replaced by Leaf (Leaf x) if x is to be kept, or Empty if it is not.

3.3 All Pairs

Task 3.3 (5 pts). Define a function

```
allpairs : 'a tree * 'b tree -> ('a * 'b) tree
```

such that allpairs(t1, t2) contains the pair (x,y) for every element x of t1 and y of t2. The order of the pairs is unspecified.

For example,

```
allpairs (Node(Leaf 1, Leaf 2), Node(Leaf "a", Leaf "b"));
== Node (Node (Leaf (1,"a"), Leaf (1,"b")), Node (Leaf (2,"a"), Leaf (2,"b")))
```

3.4 Partnr

You are writing an app to help students find study partners. Each student fills out a survey, which, for simplicity, we will assume consists of the four questions in Figure 1. That is, each person provides a tuple

```
(username, answer1, answer2, answer3, answer4) : string * int * int * int * int
```

where username is a string identifying the person, and each answerN is the number of that person's answer to question N.

For convenience, we abbreviate the tuple of int's by the type abbreviation

```
type answers = int * int * int * int
```

For example,

```
("drl", 5, 2, 1, 2) : string * answers
```

means that my answers are "all of the above", "in the computer lab", "music", and "let's do it live."

Scoring Functions A *scoring function* is a function that takes two people's answers and computes a score, which is a floating point number, where higher numbers indicate higher study partner compatibility.

For example, here is a simple scoring function, which totals up the number of answers that two people have in common:

```
fun same(x : int, y : int) : real =
   case x = y of
        true => 1.0
        | false => 0.0

fun count_same ((a1,a2,a3,a4) : answers , (a1',a2',a3',a4') : answers) : real =
        same (a1,a1') + same (a2,a2') + same (a3,a3') + same (a4,a4')
```

Scoring functions should be symmetric: score(a1,a2) == score(a2,a1).

Task 3.4 (2 pts). Write another scoring function my_scoring, which implements some alternative compatibility scoring of your choice. Explain in a comment why you think your scoring mechanism would provide good results. For example, other scoring functions might give different weights to different questions, or allow "fuzzy matching" of answers (e.g. if one person likes to study in the morning and another likes to study any time, maybe they are somewhat compatible.)

Matching Your goal is to analyze the given data and report a ranked list of possible study partners. You are given a scoring function and a cutoff, and should only report possible partners whose scores are above the cutoff.

Task 3.5 (13 pts). Write a function

where similarity is a scoring function, cutoff is a real number, people is the input data for all of the users. matches (similarity, cutoff, people) should compute a tree of pairs (person1,person2,score) where

- each score is the similarity score of person1 and person2
- the tree is sorted from highest scores to lowest scores
- only pairs of people whose score is bigger than cutoff are included
- the tree never contains a pair of people of the form (person1, person1,) or both the pair (person1, person2,) and the pair (person2, person1,).

Hints:

• Start by making the tree of all pairs of people.

- 1. What time do you like to do your homework?
 - 1 Morning
 - 2 Afternoon
 - 3 Evening
 - 4 Late-night
 - 5 All of the above
- 2. What's your preferred brand of skype?
 - 1 Hangouts
 - 2 Zoom
 - 3 Facetime
- 3. What kind of background noise helps you study?
 - 1 Music
 - 2 TV
 - 3 Quiet, please!
- 4. How do you like to spread your work for a particular class out over the week?
 - 1 Slow and steady wins the race
 - 2 Let's do it live

Figure 1: Survey

- Use < to test whether one string is less than another, according to lexicographic order.
- Use sort (described at the beginning of the handout) to sort trees.
- You can use the function **show_matches** to print the results in a nice way. For example, we have provided