

# Lecture 53 Lists

A list [int] is either

- []

"nil"  
"empty"

"cons" =  $x :: xs$ , where  $x : int$

xs is an int list

→ and that's it!

A natural number is either

- 0, or

-  $1 + k$ ,  $k \in \text{Nat}$

self-referencing

Inductive

→ and that's it!

0 nat

$0 + 1 = 1 \text{ nat}$

$0 + 1 + 1 = 2 \text{ nat}$

⋮

[ ]

: int list

(1 :: [ ]) : int list

2 :: (1 :: [ ]) : int list

~~1 :: ( )~~

2 :: 1 :: [ ]

[2, 1]

5 :: (2 :: (1 :: [ ])) : int list

5 :: 2 :: 1 :: [ ]

5 88 2 88 1 88 [ ]

[5, 2, 1]

fun  $f(l:\text{int list}) =$

case  $(l:\text{int list})$  of

[ ] =>

- - - ,  $\circ T$

: T

|  $x \circ \circ xs \Rightarrow$

- - - ,  $\circ T$

X

xs

$f(xs)$  ] Recursive  
call

$x:\text{int}$

$xs:\text{int list}$

bound variables

case [] of C]  $\Rightarrow e_1$

| x::xs  $\Rightarrow e_2$

$\mapsto e_1$

case V :: VS of C]  $\Rightarrow e_1$

| x::xs  $\Rightarrow \underline{e_2}$

$\mapsto$

$e_2$  with V for x

VS for xs

(\* Purpose: count the number of numbers in a list  $x$ )

(\* E.g.  $\text{length } [5, 3, 1] = 3$  \*)

5 :: [2 :: 1 :: []]

fun length (l: int list): int =  
case l of

[] => 0

| x :: xs => 1 + length(xs)

length (2 3 0 1 3 3 ( ))

→ case 2 1 1 : ( ) of ( ) => 0 | x : ( ) => 1 + length (x)

→ 1 + length (1 1 : ( ))

→ 1 + case 1 1 : ( ) of ( ) => 0 | x : ( ) => 1 + length (x)

→ 1 + (1 + length ( ))

→ 1 + 1 + case ( ) of ( ) => 0 | --

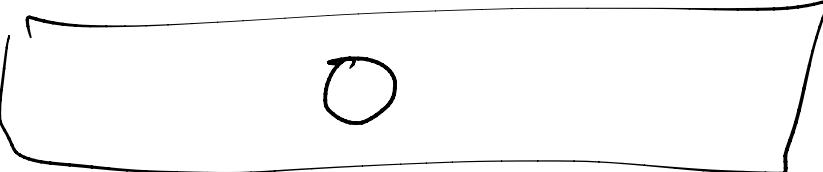
→ 1 + (1 + 0)

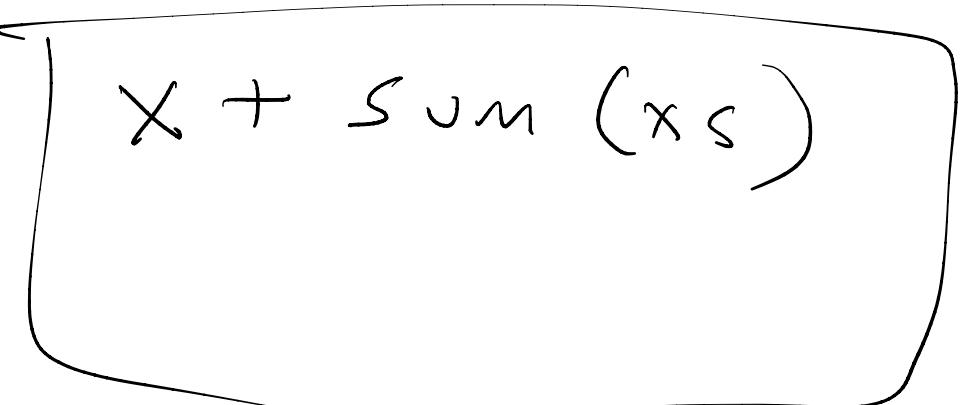
→ 2 ✓

(\* Purpose: add up all the numbers  
in a list \*)

(\* E.g.  $\text{sum} [5, 1, 2] = 8$  \*)

fun sum ( $l$ : int list): int =  
case  $l$  of

[ $\lambda \Rightarrow$  

)  $x :: xs \Rightarrow$    
 $x + \text{sum}(xs)$

Sum [5, 1, 2]

→ case 5 :: [1, 2] of C ⇒ 0 | x :: xs ⇒ x + sum(xs)

→ 5 + sum[1, 2]

→ 5 + case 1 :: [2] of C ⇒ 0 | x :: xs ⇒ x + sum(xs)

→ 5 + 1 + sum[2]

→ 5 + 1 + case 2 :: C of C ⇒ 0 | x :: xs ⇒ x + sum(xs)

→ 5 + 1 + 2 + sum[]

→ 5 + 1 + 2 + case C of C ⇒ 0 | - - -

→ 5 + (1 + (2 + 0))

∴ 8

(\* Purpose: raise everyone's salaries  
by an amount  $x$ )

(\* E.g. raiseSalaries([11, 14, 16], 4) =  
[15, 18, 20] \*)

fun raiseSalaries(l: int list, amount: int) : int list =

case l of

[] => []

| x :: xs => (x + amount) ::

raiseSalaries(xs, amount)

rs = raiseSalaries

rs([11, 14, 16], 4)

→ rs([11 :: 14 :: 16 :: []])

→ (11 + 4) :: rs([14, 16], 4)

→ 15 :: rs([14, 16], 4)

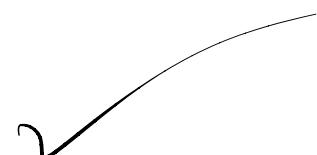
→ 15 :: (14 + 4) :: rs([16], 4)

→ 15 :: 18 :: rs([16], 4)

→ 15 :: 18 :: (16 + 4) :: rs([], 4)

→ 15 :: 18 :: 20 :: rs([], 4)

→ 15 :: 18 :: 20 :: []



for length ( $l$ : int list) : int =  
case  $l$  of  
[]  $\Rightarrow$  0

$$| \underline{x :: xs} \Rightarrow 1 + \text{length}(xs)$$

for sum ( $l$ : int list) : int =  
case  $l$  of

$$[] \Rightarrow 0$$

$$| x :: xs \Rightarrow x + \text{sum}(xs)$$

for raiseSalaries ( $l$ : int list, amount : int) : int list =

case  $l$  of

$$[] \Rightarrow []$$

$$| x :: xs \Rightarrow (x + \text{amount}) ::$$

raiseSalaries ( $xs$ ,  $\underline{\text{amount}}$ )

# Induction

"Structural induction  
for lists"

To prove Something for all  
list,

① Prove it for CJ

② For any  $x$ ,  $xS$ ,  
Assume it for a list  $xS$ ,  

---

and show it for  $x :: xS$

I H

Fusion

for all  $l, a, b$

$$rs(rs(l, a), b) \xrightarrow{q'} q_2$$

[11, 14, 16]

4

[15, 18, 20]

[17, 20, 22]

$\cong rs(l, a+b)$

For cell  $l, \underbrace{a, b}_{\text{int/list}} \quad [\text{values}]$

$$rS(rS(l, a), b) \cong rS(l, a+b)$$

Proof: by structural ind. on  $l$

Case for  $[]$ :

To show:  $rS(rS([], a), b) \cong rS([], a+b)$

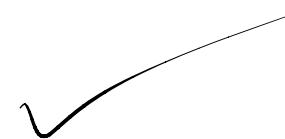
$$rS(rS([], a), b) \qquad \qquad rS([], a+b)$$

$$\mapsto rS([], b)$$

$$\mapsto []$$

$$\mapsto []$$

$$[] \cong []$$



Case for  $x :: xs$

$$\text{IH: } \text{rS}(\text{rS}(xs, a), b) \stackrel{\cong}{=} \text{rS}(xs, a+b)$$

$$\text{TS: } \text{rS}(\text{rS}(\underline{x :: xs}, a), b) \stackrel{\cong}{=} \text{rS}(x :: xs, a+b)$$

