

# Product Types / Sum Types

Types

nat  
String  
bool

Type

constructors

$\tau_1 \rightarrow \tau_2$

Gödel's  
T

int  $\rightarrow$  String  
int  $\rightarrow$  int  $\rightarrow$  int  
(int  $\rightarrow$  int)  $\rightarrow$  int

# Products

- SML tuples (1, "a", true)
- C → structs
- Python → objects

# Binary products

$\tau_1$  \*  $\tau_2$  for any types  $\tau_1$  and  $\tau_2$

(values/  
done  
exps)

$(e_1, e_2)$   
done

eager:  $e_1$  done  
 $e_2$  done

$e_1 : \tau_1 \quad e_2 : \tau_2$   
-----  
 $(e_1, e_2) : \tau_1 * \tau_2$

Int x String

(4, "hi")

(3, "a")

String x Int

("a", 7)

("b", 8)

how do you use an element  
of the product type?

→ getting out the left and right  
sides

#1  $\longrightarrow$  get the left

#2  $\longrightarrow$  get the right

#1 (4, "a")  $\longmapsto$  4

#1 (5, "b")  $\longmapsto$  5

#2 (4, "a")  $\longmapsto$  "a"

#1 ("a", 4)  $\longmapsto$  "a"

⋮  
⋮  
⋮

$\tau ::= \dots \mid \tau_1 \times \tau_2$

$e ::= \dots \mid \langle e_1, e_2 \rangle \mid \text{projleft}(e) \mid \text{projright}(e)$

SML  $(e_1, e_2) \mid \#1 e \mid \#2 e$

PFPL  $\langle e_1, e_2 \rangle \mid e.l \mid e.r$

Typing

$\Gamma \vdash e_1 : \tau_1$

$\Gamma \vdash e_2 : \tau_2$

$\Gamma \vdash \langle e_1, e_2 \rangle : \tau_1 \times \tau_2$

$\frac{\Gamma \vdash e : \tau_1 \times \tau_2}{\Gamma \vdash \text{projright}(e) : \tau_2}$

$\frac{\Gamma \vdash e : \tau_1 \times \tau_2}{\Gamma \vdash \text{projleft}(e) : \tau_1}$

$$\frac{e_1 \text{ done} \quad e_2 \text{ done}}{\langle e_1, e_2 \rangle \text{ done}}$$

Pages

$$\frac{e_1 \mapsto e_1'}{\langle e_1, e_2 \rangle \mapsto \langle e_1', e_2 \rangle}$$

$$\frac{e_2 \mapsto e_2' \quad e_1 \text{ done}}{\langle e_1, e_2 \rangle \mapsto \langle e_1, e_2' \rangle}$$

$$\frac{e \mapsto e'}{\text{projleft}(e) \mapsto \text{projleft}(e')}$$

$$\text{projright}(e) \mapsto \text{projright}(e')$$

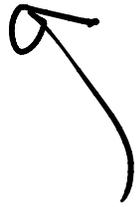
$$\frac{\langle e_1, e_2 \rangle \text{ done}}{\text{projleft}(\langle e_1, e_2 \rangle) \mapsto e_1}$$

$$\text{projright}(\langle e_1, e_2 \rangle) \mapsto e_2$$

Search

Instr

$$\frac{1+1 \mapsto 2}{\langle 1+1, 3 \rangle \mapsto \langle 2, 3 \rangle}$$

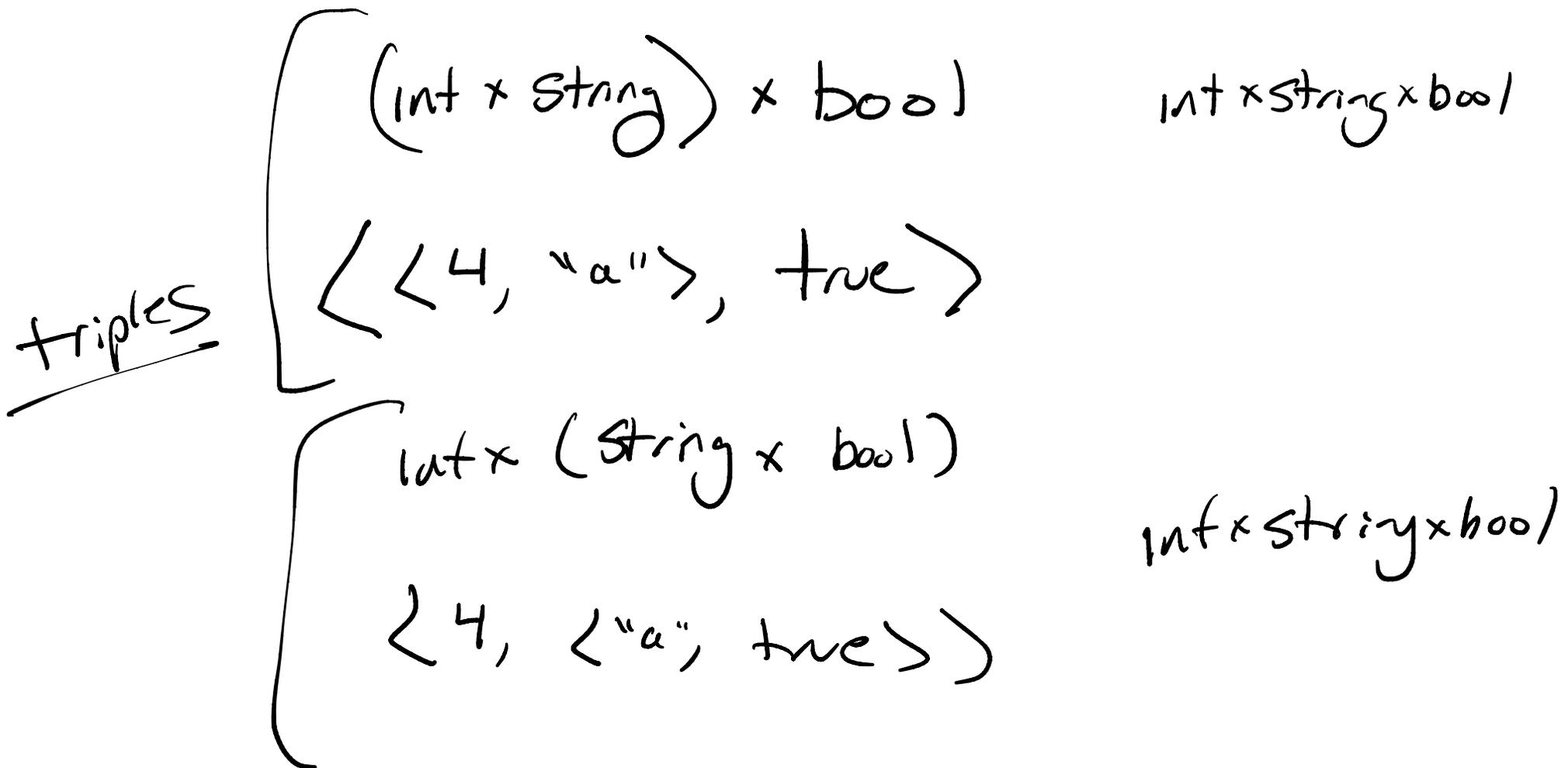


$$\text{projleft} \langle 1+1, 3 \rangle \mapsto \text{projleft} \langle 2, 3 \rangle \mapsto 2$$

$$\text{projright} \langle 1+1, 3 \rangle \mapsto \text{projright} \langle 2, 3 \rangle \mapsto 3$$

# Combining type constructors

① products and products



② products and functions

$(\text{int} \times \text{string}) \rightarrow \text{bool}$

fun f(p: int \* string): bool =

..... #1 p ..... #2 p .....

fun add(p: int \* int): int = #1 p) + (#2 p)

→ multi-argument function

fun add(x: int, y: int) = x + y

↳ add: int x int → int

---

int  $\rightarrow$  (int \* int)

multi-output  
functions

e.g. Def

$\text{fib}(1) = 1$   
 $\text{fib}(0) = 1$   
 $\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$

fun fastfib(n: int): int \* int =

case n of

0  $\Rightarrow$  (0, 1)

1  $\Rightarrow$  let val p = fastfib(n-1)

in

(#2 p, (#1 p) + (#2 p))

end

Nullary product  
→ 0 types

binary  
 $\tau_1 \times \tau_2 \mid \langle e_1, e_2 \rangle$

$\tau ::= \dots \mid \text{unit}$   
 $e ::= \dots \mid \langle \rangle$

---

$\Gamma \vdash \langle \rangle : \text{unit}$

---

$\langle \rangle \text{ done}$

Labeled products  $\rightarrow$  you choose the labels!

$e: \tau_1 \times \tau_2$  labels are generic:  
#1  $e$  #2  $e$   
projleft( $e$ ) projright( $e$ )  
 $e.l$   $e.r$

binary

$e: \{ x:int, y:int, z:bool \}$   
 $e = \{ x=4, y=7, z=true \}$   
 $\begin{matrix} e.x \\ e.y \\ e.z \end{matrix}$

Struct / object / module

{ x: int,

y: String

print: int → string

} fields

}

Products

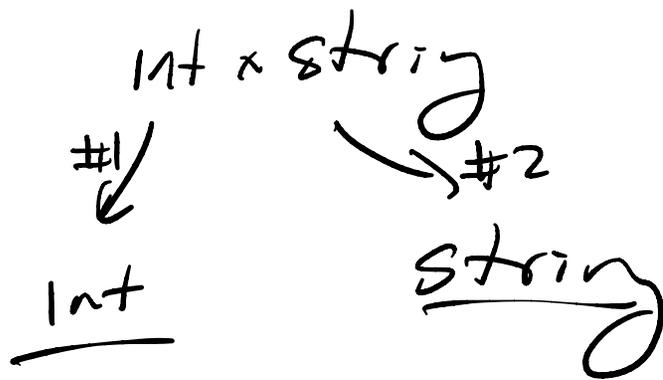
VS

Lists

fixed  
length

but

heterogeneous



arbitrary  
length  
but

homogeneous

[1, 2, 3, 7, ...]

process w/  
a loop  
where x: int  
ranges over each

# Sum Types

→ a type that  
is either  
one kind of thing  
or another

A natural number  
is either

- 0
- $1 + n$ , where...

A list  
is either

- []
- $x :: xs$  where

A `TypeError` is either

[HW2/3]

- WellTyped(t) where  $t: \text{typ}$

or - IllTyped

---

$\text{TypeError} := \underline{\text{typ}} + \underline{\text{Unit}}$

→  $\text{WellTyped}(t) := \text{Inleft}(t)$

→  $\text{IllTyped}$  :=  $\text{Inright}(\underline{\langle \rangle})$

A result is either

- Done, or

- Stopped( $e'$ ) where  $e' : \text{exp}$

---

result := unit + exp

→ Done := inleft <>

→ Stopped( $e'$ ) := inright( $e'$ )

A `typOrErr` [HW1] is either

- WellTyped ( $t$ ), where  $t: \text{typ}$

- IllTyped ( $s$ ), where  $s: \text{String}$

$\text{typOrErr} := \underline{\text{typ}} + \underline{\text{String}}$

$\text{WellTyped}(t) := \text{inleft}(t)$

$\text{IllTyped}(s) := \text{inright}(s)$

case ( $e$ ) of  
 $\{ \text{inleft}(x) \Rightarrow \dots, \text{inright}(y) \Rightarrow \dots \}$

A boolean is either

- true,

or - false

bool := unit + unit

→ true := inleft  $\langle \rangle$

→ false := inright  $\langle \rangle$

if b  
then . e1 [  $\langle \rangle$  / x ]  
else . e2 [  $\langle \rangle$  / y ]

case (b: bool) of  
{ inleft x  $\Rightarrow$  . e1  
inright y  $\Rightarrow$  . e2 }  
x, y: unit

type  $\tau ::= \dots \mid \tau_1 + \tau_2$

"either a  $\tau_1$ , or a  $\tau_2$ ,  
but tagged with  
a constructor"

exp  $e ::= \dots \mid \text{inleft}(e) \mid \text{inright}(e)$

$\mid \text{case}(e) \{ x.e_1, y.e_2 \}$  } abstract

or

case e of

{ inleft x  $\Rightarrow$  e<sub>1</sub>

inright y  $\Rightarrow$  e<sub>2</sub> }

concrete

# Typing

$$\frac{\Gamma \vdash e : \tau_1}{\Gamma \vdash \underline{\text{inleft}}(e) : \tau_1 + \tau_2}$$

$$\frac{\Gamma \vdash e : \tau_2}{\Gamma \vdash \underline{\text{inright}}(e) : \tau_1 + \tau_2}$$

"generic constructor names"

$$\frac{\Gamma \vdash e : \tau_1 + \tau_2 \quad \Gamma, x : \tau_1 \vdash e_1 : \tau \quad \Gamma, y : \tau_2 \vdash e_2 : \tau}{\Gamma \vdash \text{case}(e) \{ x.e_1 \mid y.e_2 \} : \tau}$$

case e of

{ inleft x  $\Rightarrow$  e<sub>1</sub>

inright y  $\Rightarrow$  e<sub>2</sub> }

$$\boxed{\text{Op Sem}} \xrightarrow[\text{inleft}(e) \text{ done}]{e \text{ done}}$$

$$\xrightarrow[\text{inright}(e) \text{ done}]{e \text{ done}}$$

$$\frac{e \mapsto e'}{\text{inleft}(e) \mapsto \text{inleft}(e')}$$

$$\frac{e \mapsto e'}{\text{inright}(e) \mapsto \text{inright}(e')}$$

search

$$\frac{e \mapsto e'}{\text{case}(e) \{x.e_1, y.e_2\} \mapsto \text{case}(e') \{x.e_1, y.e_2\}}$$

instr

$$\frac{e \text{ done}}{\text{case}(\text{inleft } e) \{x.e_1, y.e_2\} \mapsto e_1[e/x]}$$

$$\text{case}(\text{inright } e) \{x.e_1, y.e_2\} \mapsto e_2[e/x]$$

# Simulate sum types w/ products

int + string

want

case e of  
 inleft(x: int) => e<sub>1</sub>  
 | inright(y: string) => e<sub>2</sub>

Σ tag: bool,  
 S: { x: int  
     y: string }

inl = true  
 inr = false

if x  
 true e<sub>1</sub>[s.x] (s.y)  
 else e<sub>2</sub>[s.y] [s.x]

inleft(e) = { t = true, x = e, y = junk }  
 inright(e) = { t = false, x = junk, y = e }